

# 1-means clustering and conductance

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Network community detection with conductance

The relation to  $k$ -means clustering

Algorithms

Experiments

Conclusions

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## **Global community detection**

Given a network, find all tightly connected sets of nodes (communities).

## **Local community detection**

Given a network and a seed node, find the community/communities containing that seed. *Without inspecting the whole graph.*



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## Graphs

$$G = (V, E),$$

$$a_{ij} = a_{ji} = 1 \text{ if } (i, j) \in E \text{ else } 0$$

## Score function

$$\phi_G : \mathcal{C}(G) \rightarrow \mathbb{R}$$

Note: I'll consider sets and vectors interchangeably, so  $\mathcal{C}(G) = \mathcal{P}(V)$  or  $\mathcal{C}(G) = \mathbb{R}^V$ .



## Definition

Fraction of incident edges leaving the community

$$\phi(c) = \frac{\#\{(i,j) \in E \mid i \in c, j \notin c\}}{\#\{(i,j) \in E \mid i \in c, j \in V\}},$$

or

$$\phi(c) = 1 - \frac{\sum_{i,j \in V} c_i a_{ij} c_j}{\sum_{i,j \in V} c_i a_{ij}}$$

where  $c_i \in \{0, 1\}$ .

Very popular objective for finding network communities.



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## As an optimization problem

$$\underset{c}{\text{minimize}} \quad \phi(c)$$

$$\text{subject to} \quad c_i \in \{0, 1\} \quad \text{for all } i.$$

## Karush-Kuhn-Tucker conditions

$c$  is a local optimum if for all  $c_i$

$$0 \leq c_i \leq 1, \text{ and}$$

$$\frac{\partial \phi(c)}{\partial c_i} \leq 0 \quad \text{if } c_i < 1, \text{ and}$$

$$\frac{\partial \phi(c)}{\partial c_i} \geq 0 \quad \text{if } c_i > 0.$$



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$c$  is a local optimum if for all  $c_i$

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$$\nabla \phi(c)_i \leq 0 \quad \text{if } c_i = 1.$$

## Local optima are discrete

If  $c$  as a strict local minimum of  $\phi$ , then  $c_i \in \{0, 1\}$  for all  $i$ .

### Proof sketch

Look at  $\phi$  as a function of a single  $c_i$ :

$$\phi(c_i) = \frac{\alpha_1 + \alpha_2 c_i + \alpha_3 c_i^2}{\alpha_4 + \alpha_5 c_i}.$$

If  $0 < c_i < 1$  and  $\phi'(c_i) = 0$ , then  $\phi''(c_i) = 2\alpha_3/(\alpha_4 + \alpha_5 c_i)^3 \geq 0$ .

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## k-means clustering

$$\underset{c}{\text{minimize}} \quad \sum_{i=1}^n \sum_{j=1}^k c_{ij} \|x_i - \mu_j\|_2^2$$

Subject to the constraint that exactly one  $c_{ij}$  is 1 for every  $i$ .

## 1-means clustering

$$\underset{c}{\text{minimize}} \quad \sum_i w_i (c_i \|x_i - \mu\|_2^2 + (1 - c_i) \lambda_i)$$



## weighted $k$ -means clustering

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## 1-means clustering

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## Optimal $\mu$

Fix cluster assignment  $c_i$ , then

$$\mu = \frac{\sum_i w_i c_i x_i}{\sum_i w_i c_i}.$$

## Optimal $c$

Fix  $\mu$ , then  $c_i$  is 1 if  $\|x_i - \mu\| < \lambda_i$ , and 0 otherwise.

## Kernels

$$K(i, j) = \langle x_i, x_j \rangle$$

$$\text{so } \|x_i - x_j\|_2^2 = K(i, i) + K(j, j) - 2K(i, j).$$

## Implicit centroid

The centroid is then a linear combination of points,  $\mu = \sum_i \mu_i x_i$ , giving

$$\|x_i - \mu\|_2^2 = K(i, i) - 2 \sum_j \mu_j K(i, j) + \sum_{j, k} \mu_j K(j, k) \mu_k.$$

Optimal  $\mu$  becomes

$$\mu_i = \frac{w_i c_i}{\sum_j w_j c_j}.$$

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$$\underset{c}{\text{minimize}} \quad \sum_i w_i (c_i \|x_i - \mu\|_2^2 + (1 - c_i) \lambda_i)$$



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## 1-means objective

$$\begin{aligned} \text{minimize}_c \quad & \sum_i (w_i c_i K(i, i) - 2w_i c_i \sum_j \mu_j K(i, j)) \\ & + w_i c_i \sum_{j, k} \mu_j K(j, k) \mu_k + w_i (1 - c_i) \lambda_i \end{aligned}$$



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## 1-means objective

$$\begin{aligned} \underset{c}{\text{minimize}} \quad & \sum_i w_i c_i (K(i, i) - \lambda_i) + \sum_i w_i \lambda_i \\ & - \frac{\sum_{i, j} w_i c_i w_j c_j K(i, j)}{\sum_i w_i c_i}. \end{aligned}$$



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## 1-means objective

$$\underset{c}{\text{minimize}} \quad 1 - \frac{\sum_{i, j} w_i c_i w_j c_j K(i, j)}{\sum_i w_i c_i},$$

taking  $\lambda_i = K(i, i)$ .



## Idea

$$K = W^{-1}AW^{-1},$$

$$w_i = \sum_j a_{ij}$$

turns the objective into

$$\underset{c}{\text{minimize}} \quad 1 - \frac{\sum_{i,j} c_i c_j a_{ij}}{\sum_{i,j} c_i a_{ij}} = \phi(c),$$

We get conductance!

But this kernel is not positive definite.

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## Add a diagonal

$$K = W^{-1}AW^{-1} + \sigma W^{-1}$$

The objective becomes

$$\underset{c}{\text{minimize}} \quad 1 - \frac{\sum_{i,j} c_i c_j a_{ij}}{\sum_{ij} c_i a_{ij}} - \sigma \frac{\sum_{i,j} c_i^2 a_{ij}}{\sum_{ij} c_i a_{ij}} = \phi_\sigma(c).$$

When  $c_i \in \{0, 1\}$ ,  $c_i^2 = c_i$ , so the last term is constant.



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## Relaxing the optimization problem

$$\begin{aligned} &\text{minimize} && \phi_\sigma(c) \\ &\text{subject to} && 0 \leq c_i \leq 1 \quad \text{for all } i \in V. \end{aligned} \tag{1}$$

### Theorem

When  $\sigma \geq 2$ , every discrete community  $c$  is a local optimum of (1).

### In practice

Higher  $\sigma \Rightarrow$  more clusters are local optima.



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Given some *seed nodes*  $s$  find the cluster  $c$  that contains  $s$ .

As constrained optimization

$$\begin{array}{ll} \text{minimize} & \phi_\sigma(c) \\ \text{subject to} & 0 \leq c_i \leq 1 \quad \text{for all } i \in V. \\ \text{and} & c_i = 1 \quad \text{for all } i \in s \end{array}$$

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## Gradient descent

$$c^{(0)} = s,$$
$$c^{(t+1)} = p(c^{(t)} - \alpha^{(t)} \nabla \phi_\sigma(c^{(t)})).$$

## Project onto valid set

$$p(c) = \underset{c', \text{ s.t. } 0 \leq c'_i \leq 1, c'_i \geq s_i}{\operatorname{argmin}} \|c - c'\|_2^2,$$

This is simply

$$p(c) = \max(s, \min(1, c)).$$



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## E-step:

$c_i \leftarrow 1$  if  $\|x_i - \mu\| \leq \lambda_i$ , and 0 otherwise.

## M-step:

$$\mu_i \leftarrow \frac{w_i c_i}{\sum_j w_j c_j}.$$

## Together

If you work this out, this is equivalent to

$$c^{(0)} = s,$$
$$c^{(t+1)} = \{i \mid \nabla \phi_\sigma(c^{(t)}) < 0\} \cup s.$$

This is gradient descent with infinite step size.



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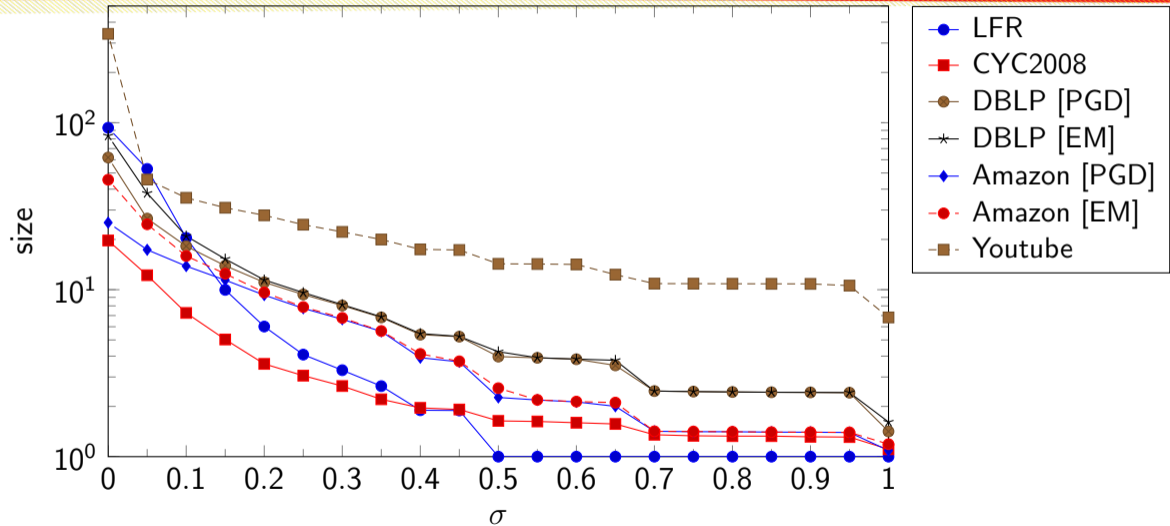
Algorithms

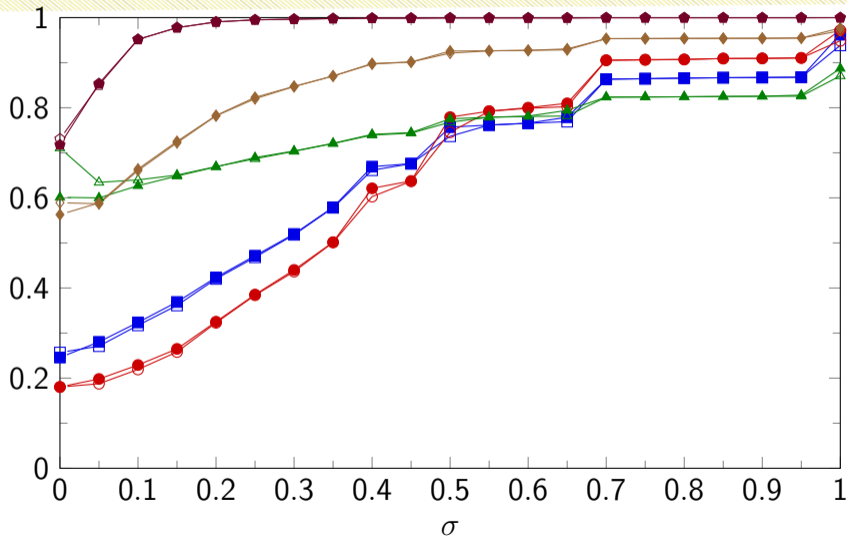
**Experiments**

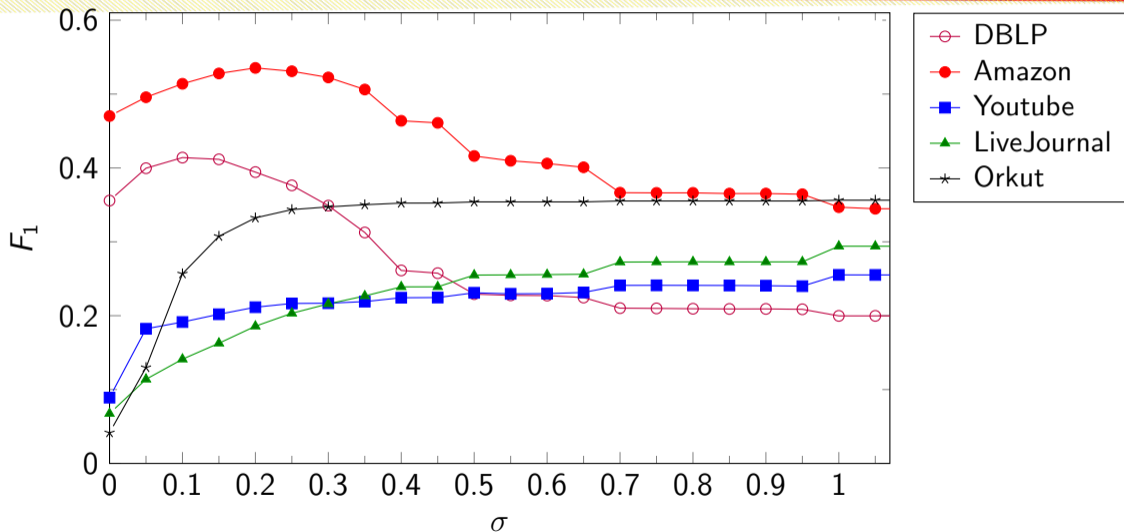
Conclusions



Dataset	#node	#edge	clus.c.	#comm
LFR (om=2)	5000	25123	0.021	146
CYC2008	6230	6531	0.121	408
Amazon	334863	925872	0.079	151037
DBLP	317080	1049866	0.128	13477
Youtube	1134890	2987624	0.002	8385
LiveJournal	3997962	34681189	0.045	287512
Orkut	3072441	117185083	0.014	6288363









- Choice of  $\sigma$  is important.
- Heuristic: Pick  $\sigma$  to maximize community density.





Dataset	PGDC-0	PGDC-d	EMC-0	EMC-d	YL	HK	PPR
LFR (om=1)	<b>0.967</b>	0.185	0.868	0.187	0.203	0.040	0.041
LFR (om=2)	<b>0.483</b>	0.095	0.293	0.092	0.122	0.039	0.041
LFR (om=3)	<b>0.275</b>	0.085	0.158	0.083	0.110	0.037	0.039
LFR (om=4)	<b>0.178</b>	0.074	0.100	0.072	0.092	0.032	0.034
Karate	0.831	0.472	0.816	0.467	0.600	0.811	<b>0.914</b>
Football	<b>0.792</b>	<b>0.816</b>	0.766	<b>0.805</b>	<b>0.816</b>	0.471	0.283
Pol.Blogs	<b>0.646</b>	0.141	<b>0.661</b>	0.149	0.017	<b>0.661</b>	0.535
Pol.Books	0.596	0.187	0.622	0.197	0.225	<b>0.641</b>	<b>0.663</b>
Flickr	0.098	0.027	0.097	0.027	0.013	0.054	<b>0.118</b>
CYC	0.474	<b>0.543</b>	0.455	<b>0.543</b>	<b>0.526</b>	0.336	0.294
Amazon	0.470	<b>0.522</b>	0.425	<b>0.522</b>	0.493	0.245	0.130
DBLP	<b>0.356</b>	<b>0.369</b>	0.317	<b>0.371</b>	0.341	0.214	0.210
Youtube	0.089	<b>0.251</b>	0.073	<b>0.248</b>	0.228	0.037	0.071
LiveJournal	0.067	<b>0.262</b>	0.059	<b>0.259</b>	0.183	0.035	0.049
Orkut	0.042	<b>0.231</b>	0.033	<b>0.231</b>	0.171	0.057	0.033

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- Conductance is related to 1-means clustering.
- Similarly, EM is the limiting case of Projected Gradient Descent.
- High conductance clusters often correspond to 'true' clusters.
- ...but they can be very large.

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